

Forward–partial inverse–forward method for solving monotone inclusions: application to land use planning ¹

L. M. Briceño-Arias

Universidad Técnica Federico Santa María

ADGO 2013
Playa Blanca
15 October 2013

¹Thanks to projects MathAmSud N 13MATH01, Anillo ACT1106, FONDEF D10I1002, and FONDECYT 3120054

Problem

Problem (P)

Find $x \in \mathcal{H}$ such that $0 \in Ax + Bx + N_V x$.

Problem

Problem (P)

Find $x \in \mathcal{H}$ such that $0 \in Ax + Bx + N_Vx$.

- \mathcal{H} is a real Hilbert space,
- $A: \mathcal{H} \rightarrow 2^{\mathcal{H}}$ is maximally monotone, i.e., it is **monotone**:

$$(\forall u \in Ax)(\forall v \in Ay) \quad \langle u - v \mid x - y \rangle \geq 0$$

and its graph is **maximal** among graphs of monotone op.

- $B: \mathcal{H} \rightarrow \mathcal{H}$ is monotone and χ -lipschitzian.
- V is a closed vectorial subspace of \mathcal{H} ($N_V = V^\perp$ is the normal cone to V).

Problem

Problem (P)

Find $x \in \mathcal{H}$ such that $0 \in Ax + Bx + N_Vx$.

- \mathcal{H} is a real Hilbert space,
- $A: \mathcal{H} \rightarrow 2^{\mathcal{H}}$ is maximally monotone, i.e., it is **monotone**:

$$(\forall u \in Ax)(\forall v \in Ay) \quad \langle u - v \mid x - y \rangle \geq 0$$

and its graph is **maximal** among graphs of monotone op.

- $B: \mathcal{H} \rightarrow \mathcal{H}$ is monotone and χ -lipschitzian.
- V is a closed vectorial subspace of \mathcal{H} ($N_V = V^\perp$ is the normal cone to V).
- We suppose that the set of solutions to (P) is $Z \neq \emptyset$.

Examples

- If $A = \partial f$ and $B = \nabla g$, where f and g are convex functions, the problem (P) reduces to (+ qualification conditions)

$$\underset{x \in V}{\text{minimize}} \quad f(x) + g(x).$$

Examples

- If $A = \partial f$ and $B = \nabla g$, where f and g are convex functions, the problem (P) reduces to (+ qualification conditions)

$$\underset{x \in V}{\text{minimize}} \quad f(x) + g(x).$$

- If $\mathcal{H} = H^n$, where H is a real Hilbert space, $A = A_1 \times \cdots \times A_n$, where $A_i: H \rightarrow 2^H$ is maximally monotone, $B: (x_i)_{1 \leq i \leq n} \mapsto (Bx_i)_{1 \leq i \leq n}$, where B is single-valued, monotone, and Lipschitzian, and $V = \{x \in \mathcal{H} \mid x_1 = \cdots = x_n\}$, the problem (P) becomes

$$\text{find } x \in \mathcal{H} \quad \text{such that} \quad 0 \in \sum_{i=1}^n A_i x + Bx.$$

Particular case 1

(P) Find $x \in \mathcal{H}$ such that $0 \in Ax + Bx + N_V x$.

If $B \equiv 0$, (P) becomes

find $x \in V$ and $y \in V^\perp$ such that $y \in Ax$

Particular case 1

(P) Find $x \in \mathcal{H}$ such that $0 \in Ax + Bx + N_Vx$.

If $B \equiv 0$, (P) becomes

find $x \in V$ and $y \in V^\perp$ such that $y \in Ax$

Partial inverse of A with respect to V

$$A_V: \mathcal{H} \rightarrow 2^{\mathcal{H}}$$

$$v \in A_V u \Leftrightarrow P_V v + P_{V^\perp} u \in A(P_V u + P_{V^\perp} v)$$

$$A_{\mathcal{H}} = A \text{ and } A_{\{0\}} = A^{-1}$$

Particular case 1

(P) Find $x \in \mathcal{H}$ such that $0 \in Ax + Bx + N_Vx$.

If $B \equiv 0$, (P) becomes

find $x \in V$ and $y \in V^\perp$ such that $y \in Ax$

Partial inverse of A with respect to V

$$A_V: \mathcal{H} \rightarrow 2^{\mathcal{H}}$$

$$v \in A_V u \Leftrightarrow P_V v + P_{V^\perp} u \in A(P_V u + P_{V^\perp} v)$$

$$A_{\mathcal{H}} = A \text{ and } A_{\{0\}} = A^{-1}$$

Problem of partial inverse

find $x \in V$ and $y \in V^\perp$ such that $0 \in A_V(x + y)$.

Particular case 1

Method: Proximal point algorithm (Martinet (1970), Rockafellar (1976)) $x_{n+1} = J_{\gamma_n A_V} x_n = (\text{Id} + \gamma_n A_V)^{-1} x_n, \gamma_n > 0$

Particular case 1

Method: Proximal point algorithm (Martinet (1970), Rockafellar (1976)) $x_{n+1} = J_{\gamma_n A_V} x_n = (\text{Id} + \gamma_n A_V)^{-1} x_n, \gamma_n > 0$

Partial inverse method (Spingarn, 1983)

Let $x_0 \in V$ and $y_0 \in V^\perp$. For every $n \in \mathbb{N}$,

Step 1. Find $(p_n, q_n) \in \mathcal{H}^2$ such that $x_n + y_n = p_n + q_n$

$$\text{and } \frac{P_V q_n}{\gamma_n} + P_{V^\perp} q_n \in A\left(P_V p_n + \frac{P_{V^\perp} p_n}{\gamma_n}\right).$$

Step 2. Set $x_{n+1} = P_V p_n$ and $y_{n+1} = P_{V^\perp} q_n$. Back to step 1.

We have $x_n \rightharpoonup x$ solution to partial inverse problem.

Particular case 1

Method: Proximal point algorithm (Martinet (1970), Rockafellar (1976)) $x_{n+1} = J_{\gamma_n A_V} x_n = (\text{Id} + \gamma_n A_V)^{-1} x_n$, $\gamma_n > 0$

Partial inverse method (Spingarn, 1983)

Let $x_0 \in V$ and $y_0 \in V^\perp$. For every $n \in \mathbb{N}$,

Step 1. Find $(p_n, q_n) \in \mathcal{H}^2$ such that $x_n + y_n = p_n + q_n$

$$\text{and } \frac{P_V q_n}{\gamma_n} + P_{V^\perp} q_n \in A\left(P_V p_n + \frac{P_{V^\perp} p_n}{\gamma_n}\right).$$

Step 2. Set $x_{n+1} = P_V p_n$ and $y_{n+1} = P_{V^\perp} q_n$. Back to step 1.

We have $x_n \rightarrow x$ solution to partial inverse problem.

$$\text{If } \gamma_n \equiv 1, \quad \begin{cases} x_{n+1} = P_V J_A(x_n + y_n) \\ y_{n+1} = P_{V^\perp}(x_n + y_n - J_A(x_n + y_n)). \end{cases}$$

Particular case 2

(P) Find $x \in \mathcal{H}$ such that $0 \in Ax + Bx + N_V x$.

Particular case 2

(P) Find $x \in \mathcal{H}$ such that $0 \in Ax + Bx + N_V x$.

If $V = \mathcal{H}$, (P) becomes

(P2) Find $x \in \mathcal{H}$ such that $0 \in Ax + Bx$.

Particular case 2

(P) Find $x \in \mathcal{H}$ such that $0 \in Ax + Bx + N_V x$.

If $V = \mathcal{H}$, (P) becomes

(P2) Find $x \in \mathcal{H}$ such that $0 \in Ax + Bx$.

Method: Forward-backward-forward (Tseng (2000)):

$$(\forall n \in \mathbb{N}) \quad \left\{ \begin{array}{l} \gamma_n \in]0, \chi^{-1}[\\ y_n = z_n - \gamma_n Bz_n \\ p_n = J_{\gamma_n A} y_n \\ q_n = p_n - \gamma_n Bp_n \\ z_{n+1} = z_n - y_n + q_n. \end{array} \right.$$

Particular case 2

$$(P) \quad \text{Find } x \in \mathcal{H} \quad \text{such that} \quad 0 \in Ax + Bx + N_V x.$$

If $V = \mathcal{H}$, (P) becomes

$$(P2) \quad \text{Find } x \in \mathcal{H} \quad \text{such that} \quad 0 \in Ax + Bx.$$

Method: Forward-backward-forward (Tseng (2000)):

$$(\forall n \in \mathbb{N}) \quad \left\{ \begin{array}{l} \gamma_n \in]0, \chi^{-1}[\\ y_n = z_n - \gamma_n Bz_n \\ p_n = J_{\gamma_n A} y_n \\ q_n = p_n - \gamma_n Bp_n \\ z_{n+1} = z_n - y_n + q_n. \end{array} \right.$$

We obtain that $(x_n)_{n \in \mathbb{N}}$ converges weakly to a solution to (P2).

General case

Existing methods do not exploit the whole structure of the problem:

- Combettes (2009) and B-A & Combettes (2011) propose an algorithm that converges weakly to a solution to (P) . However, it is necessary to compute $(\text{Id} + B)^{-1}$.
- Combettes and Pesquet (2012) exploit the single-valued property of B . However, the algorithm does not take into advantage the normal cone structure and the use of product space techniques generates several additional auxiliary variables to be updated at each iteration.

Objectives

- To propose a new convergent method for solving problem (P) that take into advantage of all the structure of the problem.
- To generalize the previous methods: partial inverse and forward-backward-forward.
- Applications: monotone inclusions involving partial sums and land use planning and Generalized Nash equilibrium problems.

- 1 Motivation
- 2 Characterization
- 3 Algorithm and convergence
- 4 Applications

Characterization

(P) Find $x \in \mathcal{H}$ such that $0 \in Ax + Bx + N_V x$.

Characterization

(P) Find $x \in \mathcal{H}$ such that $0 \in Ax + Bx + N_Vx$.

Equivalence

x is a solution to Problem (P) if and only if

$$x \in V \text{ and } (\exists y \in V^\perp) \text{ such that}$$

$$0 \in (\lambda A)_V \underbrace{(x + \lambda(y - P_{V^\perp} Bx))}_z + \lambda P_V B P_V \underbrace{(x + \lambda(y - P_{V^\perp} Bx))}_z,$$

where $\lambda \in]0, +\infty[$.

Characterization

(P) Find $x \in \mathcal{H}$ such that $0 \in Ax + Bx + N_Vx$.

Equivalence

x is a solution to Problem (P) if and only if

$$x \in V \quad \text{and} \quad (\exists y \in V^\perp) \quad \text{such that}$$

$$0 \in (\lambda A)_V \underbrace{(x + \lambda(y - P_{V^\perp} Bx))}_Z + \lambda P_V B P_V \underbrace{(x + \lambda(y - P_{V^\perp} Bx))}_Z,$$

where $\lambda \in]0, +\infty[$.

Note that:

- 1 $(\lambda A)_V$ is maximally monotone.
- 2 $\lambda P_V B P_V$ is $\lambda\chi$ -lipschitzien and monotone.
- 3 $Z = P_V((\lambda A)_V + \lambda P_V B P_V)^{-1}(0)$.

Main result

Let $\lambda \in]0, +\infty[$, let $(\gamma_n)_{n \in \mathbb{N}}$ be a sequence in $]0, (\lambda\chi)^{-1}[$, let $x_0 \in V$ and let $y_0 \in V^\perp$. For every $n \in \mathbb{N}$,

Step 1. Find (p_n, q_n) such that $x_n - \gamma_n \lambda P_V Bx_n + \lambda y_n = p_n + \lambda q_n$

$$\text{and } \frac{P_V q_n}{\gamma_n} + P_{V^\perp} q_n \in A\left(P_V p_n + \frac{P_{V^\perp} p_n}{\gamma_n}\right).$$

Step 2. Let $x_{n+1} = P_V p_n + \gamma_n \lambda P_V (Bx_n - BP_V p_n)$

and $y_{n+1} = P_{V^\perp} q_n$. Back to Step 1.

Main result

Let $\lambda \in]0, +\infty[$, let $(\gamma_n)_{n \in \mathbb{N}}$ be a sequence in $]0, (\lambda\chi)^{-1}[$, let $x_0 \in V$ and let $y_0 \in V^\perp$. For every $n \in \mathbb{N}$,

Step 1. Find (p_n, q_n) such that $x_n - \gamma_n \lambda P_V Bx_n + \lambda y_n = p_n + \lambda q_n$

$$\text{and } \frac{P_V q_n}{\gamma_n} + P_{V^\perp} q_n \in A \left(P_V p_n + \frac{P_{V^\perp} p_n}{\gamma_n} \right).$$

Step 2. Let $x_{n+1} = P_V p_n + \gamma_n \lambda P_V (Bx_n - BP_V p_n)$

and $y_{n+1} = P_{V^\perp} q_n$. Back to Step 1.

Then, the sequences $(x_n)_{n \in \mathbb{N}}$ and $(y_n)_{n \in \mathbb{N}}$ are in V and V^\perp , respectively, and for a solution $\bar{x} \in Z$ and $\bar{y} \in V^\perp \cap (A\bar{x} + P_V B\bar{x})$ we have:

- 1 $x_n \rightarrow \bar{x}$ and $y_n \rightarrow \bar{y}$.
- 2 $x_{n+1} - x_n \rightarrow 0$ and $y_{n+1} - y_n \rightarrow 0$.
- 3 $P_V Bx_n \rightarrow P_V B\bar{x}$.

Particular cases

- If $\lambda = 1$ and $B \equiv 0$, the proposed method becomes the partial inverse algorithm (Spingarn, 1983) that solves
find $x \in V$ and $y \in V^\perp$ such that $y \in Ax$.

Particular cases

- If $\lambda = 1$ and $B \equiv 0$, the proposed method becomes the partial inverse algorithm (Spingarn, 1983) that solves

$$\text{find } x \in V \text{ and } y \in V^\perp \text{ such that } y \in Ax.$$
- Step 1 is not always easy to compute. If we set $\gamma_n \equiv 1$ we obtain

Let $\lambda \in]0, \beta[$, $x_0 \in V$, $y_0 \in V^\perp$.

$$(\forall n \in \mathbb{N}) \quad \left[\begin{array}{l} s_n = x_n - \lambda P_V B x_n + \lambda y_n \\ p_n = J_{\lambda A} s_n \\ r_n = p_n - \lambda P_V B P_V p_n \\ q_n = (s_n - p_n) / \lambda \\ x_{n+1} = P_V (x_n - s_n + r_n) \\ y_{n+1} = P_{V^\perp} q_n. \end{array} \right.$$



Particular cases

- If $\lambda = 1$ and $B \equiv 0$, the proposed method becomes the partial inverse algorithm (Spingarn, 1983) that solves

$$\text{find } x \in V \text{ and } y \in V^\perp \text{ such that } y \in Ax.$$
- Step 1 is not always easy to compute. If we set $\gamma_n \equiv 1$ we obtain

Let $\lambda \in]0, \beta[$, $x_0 \in V$, $y_0 \in V^\perp$.

$$(\forall n \in \mathbb{N}) \quad \left[\begin{array}{l} s_n = x_n - \lambda P_V B x_n + \lambda y_n \\ p_n = J_{\lambda A} s_n \\ r_n = p_n - \lambda P_V B P_V p_n \\ q_n = (s_n - p_n) / \lambda \\ x_{n+1} = P_V (x_n - s_n + r_n) \\ y_{n+1} = P_{V^\perp} q_n. \end{array} \right.$$

- If $V = \mathcal{H}$ we obtain the forward–backward–forward splitting (Tseng, 2000) with a constant step size.

- 1 Motivation
- 2 Characterization
- 3 Algorithm and convergence
- 4 Applications

Composite monotone inclusions involving partial sums

Definition

Let $A: \mathcal{H} \rightarrow 2^{\mathcal{H}}$ and $B: \mathcal{H} \rightarrow 2^{\mathcal{H}}$ be maximally monotone operators and let V be a closed vectorial subspace of \mathcal{H} . We define the **partial sum of A and B with respect to V** by

$$A \square_V B = (A_V + B_V)_V.$$

Note that

- 1 $A \square_{\mathcal{H}} B = A + B$
- 2 $A \square_{\{0\}} B = A \square B = (A^{-1} + B^{-1})^{-1}$ (parallel sum).

Monotone inclusion involving partial sums

Problem

Find $x \in H$ such that

$$z \in \mathcal{A}x + N_{\mathcal{U}}x + \sum_{i=1}^m \left(\mathcal{L}_i^* (\mathcal{B}_i \square_{\mathcal{V}_i^\perp} \mathcal{D}_i) (\mathcal{L}_i x - r_i) + \mathcal{L}_i^* N_{\mathcal{V}_i} (\mathcal{L}_i x - r_i) \right) + \mathcal{C}x$$

For every $i \in \{1, \dots, m\}$:

- 1 \mathcal{U} and \mathcal{V}_i are closed vectorial subspaces of real Hilbert spaces H and G_i , respectively.
- 2 $\mathcal{A}: H \rightarrow 2^H$ and $\mathcal{B}_i: G_i \rightarrow 2^{G_i}$ are maximally monotone.
- 3 $\mathcal{D}_i: G_i \rightarrow 2^{G_i}$ is monotone and $(\mathcal{D}_i)_{\mathcal{V}_i^\perp}$ is ν_i -lipschitzian.
- 4 $\mathcal{C}: H \rightarrow H$ is monotone and μ -lipschitzian.
- 5 $\mathcal{L}_i: H \rightarrow G_i$ is linear and bounded.
- 6 $z \in H$ and $r_i \in G_i$.

Particular case

If $\mathcal{U} = H$ and, for every $i \in \{1, \dots, m\}$, $\mathcal{V}_i = \mathbf{G}_i$ the problem reduces to

Find $x \in H$ such that

$$z \in \mathcal{A}x + \sum_{i=1}^m \left(\mathcal{L}_i^*(\mathcal{B}_i \square \mathcal{D}_i)(\mathcal{L}_i x - r_i) \right) + \mathcal{C}x,$$

which is solved in Combettes-Pesquet (2012).

Equivalent formulation

KKT conditions yields

Primal-Dual inclusion

$$\begin{bmatrix} 0 \\ 0 \\ \vdots \\ 0 \end{bmatrix} \in \underbrace{\begin{bmatrix} A & 0 & \dots & 0 \\ 0 & (\mathcal{B}_1)_{\mathcal{V}_1^\perp} & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \dots & (\mathcal{B}_m)_{\mathcal{V}_m^\perp} \end{bmatrix}}_A \begin{bmatrix} x \\ v_1 \\ \vdots \\ v_m \end{bmatrix} + \underbrace{\begin{bmatrix} \mathcal{C} & \mathcal{L}_1^* & \dots & \mathcal{L}_m^* \\ -\mathcal{L}_1 & (\mathcal{D}_1)_{\mathcal{V}_1^\perp} & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ -\mathcal{L}_m & 0 & \dots & (\mathcal{D}_m)_{\mathcal{V}_m^\perp} \end{bmatrix}}_B \begin{bmatrix} x \\ v_1 \\ \vdots \\ v_m \end{bmatrix} + \underbrace{\begin{bmatrix} N_{\mathcal{U}} & 0 & \dots & 0 \\ 0 & N_{\mathcal{V}_1} & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \dots & N_{\mathcal{V}_m} \end{bmatrix}}_{N_W} \begin{bmatrix} x \\ v_1 \\ \vdots \\ v_m \end{bmatrix}$$

Equivalent formulation

KKT conditions yields

Primal-Dual inclusion

$$\text{Find } z \in \mathcal{H} \quad \text{s.t.} \quad 0 \in Az + Bz + N_W z,$$

where

- $\mathcal{H} = H \times G_1 \times \cdots \times G_m$.
- $A: \mathcal{H} \rightarrow 2^{\mathcal{H}}$ is maximally monotone.
- $B: \mathcal{H} \rightarrow \mathcal{H}$ is monotone and lipschitzian.
- $W = \mathcal{U} \times \mathcal{V}_1 \times \cdots \times \mathcal{V}_m$.

Equivalent formulation

KKT conditions yields

Primal-Dual inclusion

$$\text{Find } z \in \mathcal{H} \quad \text{s.t.} \quad 0 \in Az + Bz + N_W z,$$

where

- $\mathcal{H} = H \times G_1 \times \cdots \times G_m$.
- $A: \mathcal{H} \rightarrow 2^{\mathcal{H}}$ is maximally monotone.
- $B: \mathcal{H} \rightarrow \mathcal{H}$ is monotone and lipschitzian.
- $W = \mathcal{U} \times \mathcal{V}_1 \times \cdots \times \mathcal{V}_m$.

Our algorithm applied in this case gives a splitting convergent algorithm.

Land use planning: Notation

- C : finite set of types of households (or agents). For every $h \in C$, $H_h > 0$ is the number of agents of type h .
- N : finite set of zones. For every $i \in N$, $S_i > 0$ is the number of available houses in the zone i .
- x_{hi} : % of agents of type h which will be localized in the zone i (variable to obtain).
- **Constraints**: for every $i \in N$ and $h \in C$, $\sum_{h \in C} H_h x_{hi} \leq S_i$ and $\sum_{i \in N} x_{hi} = 1$.
- z_{hi} : perceived utility of agents type h with respect to the zone i . It depends on the localization of the agents of all types (externality):

$$z_{hi} = \gamma_{hi} + \sum_{g \in C} \alpha_{hg} x_{gi}.$$

Land use planning: Notation

|C|-players game: we denote $x^h = (x_{hi})_{i \in N}$ and x^{-h} as usual.

- Strategy set: $\Delta = \{x \in [0, +\infty[^{|\mathcal{N}|} \mid \sum_{i \in N} x_i = 1\}$
- Payoff agent type h :

$$\begin{aligned} F_h(x^h, x^{-h}) &= \sum_{i \in N} (x_{hi} z_{hi} - \mu x_{hi} (\ln(x_{hi}) - 1)) \\ &= \sum_{i \in N} (x_{hi} \gamma_{hi} - \mu x_{hi} (\ln(x_{hi}) - 1) + \alpha_{hh} x_{hi}^2) \\ &\quad + \sum_{i \in N} \sum_{g \neq h} \alpha_{hg} x_{hi} x_{gi} \end{aligned}$$

- Shared constraints: For every $i \in N$,
 $X_i = \{x = (x^h)_{h \in C} \in [0, +\infty[^{|\mathcal{C}|+|N|} \mid \sum_{h \in C} H_h x_{hi} \leq S_i\}$.
 Denote by $X = \bigcap_{i \in N} X_i$.

Land use planning: Model

Generalized Nash equilibrium (GNE)

Find $x = (x^h)_{h \in C} \in \Delta^{|M|} \cap X$ such that

$$(\forall h \in C)(\forall y^h \in \Delta \cap X_{x^{-h}}) \quad F_h(x^h, x^{-h}) \geq F_h(y^h, x^{-h}).$$

Land use planning: Model

Generalized Nash equilibrium (GNE)

Find $x = (x^h)_{h \in C} \in \Delta^{|N|} \cap X$ such that

$$(\forall h \in C)(\forall y^h \in \Delta \cap X_{x^{-h}}) \quad F_h(x^h, x^{-h}) \geq F_h(y^h, x^{-h}).$$

Particular case: If, for every g, h , $\alpha_{hg} = 0$, our game becomes the potential game:

$$\underset{x \in X \cap \Delta^{|N|}}{\text{maximize}} \quad \sum_{h \in C} \sum_{i \in N} (x_{hi} \gamma_{hi} - \mu x_{hi} (\ln(x_{hi}) - 1)),$$

which has been proposed by Wilson (1967) (see also Roy (2004)).

Land use planning: Model

Generalized Nash equilibrium (GNE)

Find $x = (x^h)_{h \in C} \in \Delta^{|N|} \cap X$ such that

$$(\forall h \in C)(\forall y^h \in \Delta \cap X_{x^{-h}}) \quad F_h(x^h, x^{-h}) \geq F_h(y^h, x^{-h}).$$

Particular case: If, for every g, h , $\alpha_{hg} = 0$, our game becomes the potential game:

$$\underset{x \in X \cap \Delta^{|N|}}{\text{maximize}} \quad \sum_{h \in C} \sum_{i \in N} (x_{hi} \gamma_{hi} - \mu x_{hi} (\ln(x_{hi}) - 1)),$$

which has been proposed by Wilson (1967) (see also Roy (2004)).

Existing methods for solving (GNE) use lagrangian multipliers or penalty methods (see Facchinei-Fischer-Piccialli, 2009; Facchinei-Kansow, 2007; Pang-Fukushima, 2005, ...)

Land use planning: Equivalent formulation

Define

$$\Phi: \mathbb{R}^{|C|+|N|} \rightarrow \mathbb{R}^{|C|+|N|}: \mathbf{x} = (\mathbf{x}^h)_{h \in C} \mapsto (-\nabla_{\mathbf{x}^h} F_h(\mathbf{x}^h, \mathbf{x}^{-h}))_{h \in C}.$$

Land use planning: Equivalent formulation

Define

$$\Phi: \mathbb{R}^{|\mathcal{C}|+|N|} \rightarrow \mathbb{R}^{|\mathcal{C}|+|N|}: \mathbf{x} = (\mathbf{x}^h)_{h \in \mathcal{C}} \mapsto (-\nabla_{\mathbf{x}^h} F_h(\mathbf{x}^h, \mathbf{x}^{-h}))_{h \in \mathcal{C}}.$$

Since, for every $h \in \mathcal{C}$ and \mathbf{x}^{-h} , $-F_h(\cdot, \mathbf{x}^{-h})$ is convex, it is enough to solve:

$$\begin{aligned} \text{find } \mathbf{x} = (\mathbf{x}^h)_{h \in \mathcal{C}} \in X \cap \Delta^{|N|} \quad \text{such that} \\ (\forall \mathbf{y} \in X \cap \Delta^{|N|}) \quad \langle \Phi(\mathbf{x}) \mid \mathbf{y} - \mathbf{x} \rangle \geq 0, \end{aligned}$$

or, equivalently, (qualification conditions hold)

Monotone inclusion

$$\text{find } \mathbf{x} \in X \quad \text{such that} \quad 0 \in \Phi(\mathbf{x}) + N_{\Delta^{|N|}}(\mathbf{x}) + N_X(\mathbf{x}).$$

Land use planning: Equivalent formulation

Define

$$\Phi: \mathbb{R}^{|\mathcal{C}|+|\mathcal{N}|} \rightarrow \mathbb{R}^{|\mathcal{C}|+|\mathcal{N}|}: \mathbf{x} = (x^h)_{h \in \mathcal{C}} \mapsto (-\nabla_{x^h} F_h(x^h, x^{-h}))_{h \in \mathcal{C}}.$$

Since, for every $h \in \mathcal{C}$ and x^{-h} , $-F_h(\cdot, x^{-h})$ is convex, it is enough to solve:

$$\begin{aligned} \text{find } \mathbf{x} = (x^h)_{h \in \mathcal{C}} \in X \cap \Delta^{|\mathcal{N}|} \quad \text{such that} \\ (\forall \mathbf{y} \in X \cap \Delta^{|\mathcal{N}|}) \quad \langle \Phi(\mathbf{x}) \mid \mathbf{y} - \mathbf{x} \rangle \geq 0, \end{aligned}$$

or, equivalently, (qualification conditions hold)

Monotone inclusion

$$\text{find } \mathbf{x} \in X \quad \text{such that} \quad 0 \in \Phi(\mathbf{x}) + N_{\Delta^{|\mathcal{N}|}}(\mathbf{x}) + N_X(\mathbf{x}).$$

Remark: All solution of the inclusion is NE, but not every NE is a solution of the inclusion. These special NE are called **variational equilibria** (see e.g. Facchinei-Kansow, 2007).

Land use planning: Equivalent formulation

Variable substitution: $u = x - e$, where $e = (1, \dots, 1)/|N|$.
 Define $V = \{x \in \mathbb{R}^{|\mathcal{C}|+|N|} \mid (\forall h \in \mathcal{C}) \sum_{i \in N} x_i^h = 0\}$, which is a vectorial subspace of $\mathbb{R}^{|\mathcal{C}|+|N|}$, and $\tilde{X} = X - e$. Then our inclusion becomes

Modified inclusion

find $u \in \tilde{X}$ such that $0 \in \Phi(u + e) + N_V u + N_{\tilde{X}} u$.

Land use planning: Equivalent formulation

Variable substitution: $u = x - e$, where $e = (1, \dots, 1)/|N|$. Define $V = \{x \in \mathbb{R}^{|C|+|N|} \mid (\forall h \in C) \sum_{i \in N} x_i^h = 0\}$, which is a vectorial subspace of $\mathbb{R}^{|C|+|N|}$, and $\tilde{X} = X - e$. Then our inclusion becomes

Modified inclusion

$$\text{find } u \in \tilde{X} \text{ such that } 0 \in \Phi(u + e) + N_V u + N_{\tilde{X}} u.$$

Under suitable conditions on the constants $(\alpha_{hg})_{h \in C, g \in C}$ we can guarantee that Φ is monotone. Hence, we can apply the forward-PI-forward method for finding a solution of the land use planning problem.

References

- L. M. Briceño-Arias, Forward-Douglas-Rachford splitting and forward-partial inverse method for solving monotone inclusions, <http://arxiv.org/abs/1212.5942>.
- L. M. Briceño-Arias and P. L. Combettes, A monotone+skew splitting model for composite monotone inclusions in duality, *SIAM Journal on Optimization*, vol. 21, pp. 1230–1250, 2011.
- P. L. Combettes and J.-C. Pesquet, Primal-dual splitting algorithm for solving inclusions with mixtures of composite, Lipschitzian, and parallel-sum type monotone operators, *Set-Valued and Variational Analysis*, vol. 20, no. 2, pp. 307–330, 2012.
- P. L. Combettes and J.-C. Pesquet, A proximal decomposition method for solving convex variational inverse problems, *Inverse Problems*, vol. 24, ID 065014, 27 pp., 2008.
- B. Martinet, Régularisation d'inéquations variationnelles par approximations successives, *Rev. Francaise Inf. Rech. Oper.*, pp. 154–159, 1970.
- T. Rockafellar, Monotone operators and the proximal point algorithm, *SIAM J. Control Optim.*, vol. 14, pp. 877–898, 1976.
- J. E. Spingarn, Partial inverse of a monotone operator, *Appl. Math. Optim.*, vol. 10, pp. 247–265, 1983.