Consistency of Intertemporal Decisions: Approaches through Robust and Stochastic Optimization.

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The Problem

- Optimization has been used for long to solve many relevant problems.
- But many times we face uncertainty, or data is not known exactly.
- We also face changing conditions through time.
- What happens when we use optimization models for decision making and things change with time?

The Problem

- We consider questions related to decision making in different time horizons
- One of the most typical examples is production planning decisions in different stages:
- Strategic, Tactical, Operational.
- Tactical decisions
 - Monthly planning decisions (for instance)
 - Aggregate production data and demand
 - Aggregate production decisions and processes
 - Aggregate decisions on resources and raw materials
- Operational decisions
 - Weekly (or daily) decisions for the first month (say)
 - Detailed use of resources
 - Detailed production plan

The Planning Process

- We say the plans are consistent if feasible operational decisions can be generated, considering the constraints imposed by the tactical decisions.
- But inconsistencies may appear due to:
 - Different degrees of aggregation
 - Uncertainty and variations which are not captured in the tactical plan.
- The plan: to study factors affecting the consistency of decisions and procedures to cope with this.

Outline

- A general Setting of Intertemporal Decisions
- A framework to control inconsistencies
- A practical problem: intertemporal decisions in forest management.
- A Robus Optimization and a 2-stage stochastic approach for the problem.
- Measures of sensitivity and robustness and its connections to the question

• Collaborators: Alfonso Lobos (M.Sc. student), Pamela Alvarez (Ph.D. student, UAB), Ana Batista (Ph.D. student).

- There is a tactical planning problem in T periods (months)
- At period *t*, production, resources and logistic decisions are made.
- Variables: (x_t, y_t) , x: production, y: resources
- Data parameters: ω_t .
- Cost functions: C_t .
- The tactical problem:

$$TP) \quad \begin{array}{l} \min \quad \sum\limits_{t=1}^{T} C_t(\omega_t, x_t, y_t) \\ s.t. \quad G_t(\omega_t, x_t, y_t) \leq b_t \quad t = 1, ..., T \\ H(\omega, x, y) = 0 \end{array}$$

- At the operational level, we see subperiods (weeks within the month, for instance):
- J(t): the set of subperiods in period t.
- Operational decisions: (x^o_{tk},y^o_{tk}) , $k\in J(t)$
- Operational parameters: $\bar{\omega}_{tk}$, $k \in J(t)$.
- The operational problem is affected by the tactical planning:

$$OP) \quad \begin{array}{ll} \min & \sum\limits_{k \in J(t)} \bar{C}_{tk}(\bar{\omega}_{tk}, x^o_{tk}, y^o_{tk}, x_t, y_t) \\ s.t. & \bar{G}_{tk}(\bar{\omega}_{tk}, x^o_{tk}, y^o_{tk}, x_t, y_t) \leq \bar{b}_{tk} \quad k \in J(t) \\ \bar{H}_t(\bar{\omega}_t, x^o_{tk}, y^o_{tk}, x_t, y_t) = 0 \end{array}$$

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- The short term arrives and operational planning is done.
- Can we guarantee to obtain a reasonable operational plan?
- Old question, in fact, consistency in hierarchical planning was studied initially by Bitran, Hax and Hass[1980] and several others later

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- At operational level we decide production of all three products and we "know" demand for the three.
- But at the tactical level, an aggregated planning is done, for total liters of milk to be processed over a longer horizon.
- The tactical planning define aggregated production capacity.
- However, use of production capacity depends on the detailed product, so an inconsistency might be generated between aggregated capacity and actual detailed capacity requirements.

- The solution:
- The aggregated productivity used as parameter in the tactical planning has to be computed as a weighted average of the detailed one, and the weights have to be the relative demand for the three products.
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- Nice result, but it requires exact knowledge of detailed future demand.
- How to get that information?

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• Of course, in practice we could get estimates, but they will have error...

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- For instance:
 - $s(\cdot)$ could be cost of not fulfilling operational requirements.
 - $s(\cdot)$ could be the probability of not fulfilling operational requirements.
 - $s(\cdot)$ could be a measure of stability of the operational problem: larger s means a less stable problem.

• Then, we could state the following problems:

$$\begin{array}{ll} \min & \sum_{t=1}^{T} C_t(\omega_t, x_t, y_t) \\ TP_{R1}) & s.t. & G_t(\omega_t, x_t, y_t) \leq b_t \quad t = 1, ..., T \\ & H(\omega, x, y) = 0 \\ & \quad \text{``min } s(\bar{\omega}_t, x_t, y_t)'' \end{array}$$

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• or:

$$TP_{R2}) \quad \begin{array}{ll} \min & \sum_{t=1}^{T} C_t(\omega_t, x_t, y_t) + \sum_{t=1}^{T} s(\bar{\omega}_t, x_t, y_t) \\ s.t. & G_t(\omega_t, x_t, y_t) \leq b_t \\ H(\omega, x, y) = 0 \end{array} \quad t = 1, \dots, T$$

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• These problems tries to compute tactical decisions in such a way that their impact on the operational problem is controlled.

• Consider first the format:

$$TP_{R1}) \quad \begin{array}{l} \min \quad \sum_{t=1}^{T} C_t(\omega_t, x_t, y_t) \\ s.t. \quad G_t(\omega_t, x_t, y_t) \le b_t \quad t = 1, ..., T \\ H(\omega, x, y) = 0 \\ \quad \text{``min } s(\bar{\omega}_t, x_t, y_t)'' \end{array}$$

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- To continue the explanation we introduce the specific test problem we have been using.

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- But what you get is not always what you asked for...
- We "present" now the (simplified) models, which are based on Weintraub and Epstein[2002] and others.

- Variables:
 - X_t : hours of labor for month t.
 - raw_{ct} : logs of type c to be ordered for month t.
 - r_{ct} : logs of type c processed in month t.
 - z_{mt} : inventory of lumber m in the month t.
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- Parameters:
 - W_t : cost of labor in month t.
 - ϕ : productivity of labor.
 - $Craw_{ct}$: cost of log type c bought in period t.
 - UX, LX: upper and lower bound on labor.
 - ME_{ct} : upper bound in the amount of logs type c the company can buy in period t.
 - h_{mt} : the storage cost of product m in the month t.
 - hw_{ct} : storage cost for log c in the month t.
 - Y_{cm} : average amount of lumber of type m obtained from a log type c.
 - d_{mt} : demand for product m the company has in the month t.

- The model cover 4 months.
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- C, set of logs types, M: set of products.
- The model seeks to determine raw material and labor need so that cost is minimized

$$\begin{split} \min & \sum_{t=1}^{4} \left(\sum_{m \in M} \left(h_{mt} z_{mt} \right) + \sum_{c \in C} \left(Craw_{ct} raw_{ct} + h_{ct} w_{ct} \right) + W_t X_t \right) \\ s.t. & LX \leq X_t \leq UX \quad \forall t \in 1, ..., 4 \\ & raw_{ct} \leq ME_{ct} \quad \forall c \in C, t \in 1, ..., 4 \\ & w_{ct} = w_{c,t-1} + raw_{ct} - r_{ct} \quad \forall c \in C, t = 2, ..., 4 \\ & z_{mt} = z_{m,t-1} + \sum_{c \in C} Y_{cm} r_{ct} - d_{mt} \quad \forall m \in M, t = 2, ..., 4 \\ & \sum_{c \in C} r_{ct} \leq \phi X_t \quad \forall t = 1, ..., 4 \\ & z_{mt} \geq 0, raw_{mt} \geq 0, w_{ct} \geq 0, X_t \geq 0 \end{split}$$

• Operational variables

- r'_{eci} : logs type c processed with cutting pattern e in week i.
- ex'_i : overtime at week i.
- z'_{mi} : inventory of product m in week i.
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- Operational parameters for the first four weeks:
 - α_m : percentage of acceptable shortage for product m.
 - β'_{mi} : backlog cost for product m in week i.
 - h'_{mi} : storage cost of product m in week i.
 - h'_{ci} : storage cost for log c in week i.
 - EW_i : overtime cost in week i.
 - RR'_{ci} : ctual logs of type c received in week i.
 - Y_{ecm} : yield of lumber m from logs c using cutting pattern e.
 - d_{mi} : demand for product m in i.

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- Here E is the set of detailed cutting patterns.
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- The model seeks to determine detailed operations, given resources and raw material assigned for the first month in the tactical model.

$$\begin{split} \min & \sum_{i=1}^{4} \left(EW'_{i}ex'_{i} + \sum_{m \in M} \beta'_{mi}b_{mi} \right) \\ s.t. & w'_{ci} = w'_{c,i-1} + RR'_{ci} - \sum_{e \in E_{c}} (r'_{eci}) \quad \forall c \in C, i \in 1, ..., 4 \\ & z'_{mi} = z'_{m,0} + \sum_{c \in C} \sum_{e \in E_{c}} Y_{ecm}r'_{eci} + b'_{mi} - b'_{m,i-1} - d_{mi} \quad \forall m \in M, \forall i \\ & \sum_{c \in C} \sum_{e \in E_{c}} r'_{eci} \leq \phi \frac{X_{1}}{4} \quad \forall i \\ & b'_{mi} \leq \alpha_{m}d_{mi} \quad \forall m \in M, i \in 1, ..., 4 \\ & ex'_{i} \geq 0, b_{mt} \geq 0. \end{split}$$

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- But in reality, there can be many changes and variations.
- The represented situation:
- Those proportional weekly quantities are randomly perturbed:

$$RR'_{ci} = \left(\frac{raw_{c1}}{4}\right) + \xi \ , \ \forall c \in C, i \in 1, ..., 4,$$

• where ξ is a random perturbation.

A Robust Optimization Approach

- The first idea was to use Robust Optimization.
- The aggregated yield coefficient Y_{cm} can be used to represent all the "noise" from the aggregation and variation at the operational.
- Hence, robust decisions are made at the tactical level and transferred to the operational first month.
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- Hence, robust decisions are made at the tactical level and transferred to the operational first month.
- The performance in the operational model is evaluated using Monte-Carlo simulation.
- What are robust solutions?

• Consider the following Linear Optimization problem

 $\begin{array}{ll} \min & c^T x \\ \text{s.t.} & Ax \leq b \\ & x \geq 0 \end{array},$

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- Typically we assume $A = \overline{A} + U$, $U \in D(\Gamma)$, where $D(\Gamma)$ is an "uncertainty set" parametrized in Γ .

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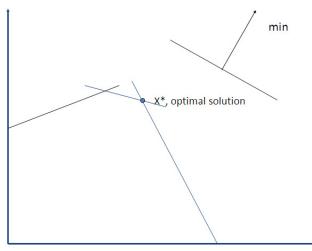
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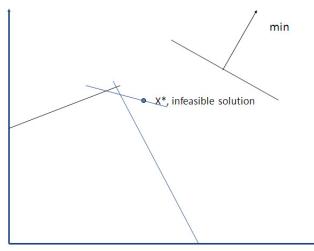
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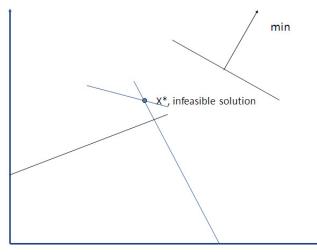
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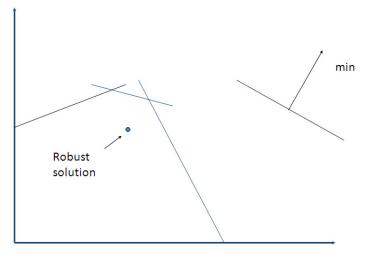
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- Γ is the degree of uncertainty (or noise) allowed.
- We look for a solution feasible for all cases of A, that is, a robust solution.









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• where, for each $i = 1, \ldots, m$,

$$\begin{array}{lll} \beta_i(x,\Gamma) &=& \max & x^T u_i \\ & \text{s.t.} & u_i \in \mathcal{D}_i(\Gamma) \end{array}$$

are the "protection functions" of the constraints.

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- This could be too conservative...
- There are other ways to vary the coefficients avoiding the worst case.
- Initial developments: Ben-Tal and Nemirovski[1998-2002], ellipsoidal uncertainty:

$$\sum_{j=1}^{n} \left(\frac{a_{ij} - \bar{a}_{ij}}{s_{ij}} \right)^2 \le \Gamma$$

Different Approaches: The "Uncertainty Budget"

• Bertsimas and Sim[2004] propose a format in which $a_{ij} \in [\bar{a}_{ij} - s_{ij}, \bar{a}_{ij} + s_{ij}]$, with the added condition that

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- If Γ es large (= n) there is no restriction on the simultaneous variation and we are in the worst case.
- For intermediate values of Γ there is an "uncertainty budget" to distribute among all coefficients as the simultaneous variation is bounded.
- With these definitions we can build the robust counterpart, which is a linear program.
- And, if there are integer variables, the robust counterpart is a mixed integer problem.

Results: Base case

- Simulation of 1,000 scenarios with different value of perturbation ξ .
- ξ is sampled from a uniform distribution, with mean $raw_{c1}/4$ and perturbed in a certain %.
- We registered when the operational problem unmet demand and when it was completely infeasible (due to processing capacity).

% perturbation	% infeasible scenarios	% of feasible scenarios with unmet demand	
20	2.60	0.00	
25	7.00	0.97	
30	14.3	6.67	
35	20.8	17.7	
40	27.8	24.4	

Results: Base case

- We now use Bertsimas and Sim robust optimization formulation for the tactical problem.
- We show the case with a variation of 30% in the log supply.
- We assume that the forest perturb with the same pattern as in the corresponding base case.
- \bullet Here are the results for some values of $\Gamma,$ the uncertainty budget in the B&S formulation, for interval widths of 15%

Value of Γ	% infeasible scenarios	% of feasible scenarios with unmet demand	
1	12.60	3.43	
3	6.80	0.00	
6	1.70	0.00	

Conclusions so far...

- A robust tactical plan increases the chances of getting feasible production at the operational level, with an increase in cost.
- From extensive simulations we could infer the right value for Γ for an acceptable feasibility level.
- However, the approach still handles both problems separated.

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- Q((raw, X), ξ) is the optimal value of the second stage, as a function of the first stage decisions (raw, X) and a random perturbation ξ.
- Hence, the problem seeks to optimize tactical decisions in such a way that the cost generated to the second stage is also taken into account.

The second stage problem

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- The second stage problem is the one we stated before.
- The random element enters in the perturbation of the disaggregated raw material supply.

$$Q((raw, X), \xi) = \min_{\substack{s.t.\\}} f_O(ex', b)$$

$$H_1(ex', b) + H_2(raw, X, ex', b) = h(\xi),$$

where the functions H_1 and H_2 represent all the operational constraints, with H_2 depending on the first stage variables.

• $h(\xi)$ represents the right-hand side as a function of the random perturbation.

A Rolling Horizon Framework

- We wanted to test the models in a rolling horizon framework
- We considered 48 months and the planning horizon moves sequentially.
- We defined certain scenarios of timber demand:



• The results we present later are for this demand pattern.

Alternative models for the 2 stages

• We considered different combinations of the tactical and operational models, which represent different views of the hierarchical decisions.

MODEL	First Stage	Second Stage	
FMA	Tactical, aggregated yield	Operational month 1	
FMD	Tactical, disaggregated yield	Operational month 1	
SMA	Simple tactical	Operational month 1 tactical inventory ag. yield month 2-4	
SMD	Simple tactical	Operational month 1 tactical inventory disag. yield	

- The disaggregated models represent a situation in which more detailed information is used on the planning.
- The second model represents the situation in which the second stage considers inventory decisions as variables that adjust to uncertainty over the whole horizon.

Solving the Problems

- We have a 2-stage Stochastic Linear Problem which we can solve in different ways.
- We used a simplified version of a real industrial problem: only a few products and cutting patterns.
- We solved the problems using an SAA (Stochastic Average Approximation) approach with 96 scenarios.
- The computations were performed in a Dell cluster with Intel E5-2470 processors with a total of 168 cores.
- The programming was developed in Python, using Gurobi for optimization.

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- 4. We solve the operational problem for that scenario and record total operational costs and solutions.

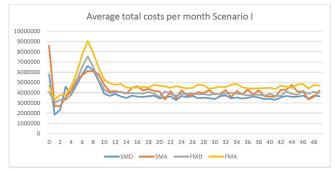
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- 5. $k \leftarrow k+1$.
- We make the operational always feasible: if there are not enough logs, we assume we purchase them on spot, with the corresponding extra cost (50% higher than the original).

Image: A matrix

Some results for demand scenario I

• Total Operational cost for the first month k for all the horizon (in monetary units)



- SMD performs better and FMA performs worst, possibly indicating that the simultaneous consideration of more information, and a full horizon in the second stage is beneficial.
- (Note: uncertainty only affects the first month).

J. Vera (PUC)

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- We see that the disaggregated model can anticipate better.
- The second model, with the inventory in the second stage, seems to perform well only under disaggregated information.
- The basic 2 stages planning model (FMA) performs well anyway.

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- What about?:
 - $s(\cdot)$ as the probability of not fulfilling operational requirements.
 - $s(\cdot)$ as a measure of stability of the operational problem: larger s means a less stable problem.
- Stability is important as a more stable operational problem will remain feasible even of some data (from the tactical) change.

Measures of Problem Stability

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- There are two kind of such measures:
 - Condition measures
 - Geometric measures
- We consider the following Linear Problem as an illustration:

$$P) \quad \begin{array}{c} z^* = \max \quad c^T x \\ \text{s.t.} \quad Ax \le b \end{array}$$

- Suppose the data is d = (A, b).
- Let $F_d := \{x \in \mathbb{R}^n : Ax \le b\} \neq \emptyset$
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- Condition number is $C(d) = \frac{\|d\|}{\rho(d)}$
- The notion was studied, for Optimization, by Renegar[1995] and further results by V[1996], Freund and Vera[1999,2003], Ordoñez and Freund[2003], and several others.
- The condition number explains, among others, sensitivity properties of the problem.

Conditioning and Problem Sensitivity

• For the optimization problem:

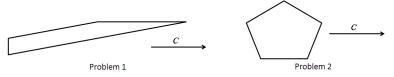
$$P) \quad \begin{array}{c} z^* = \max \quad c^T x \\ \text{s.t.} \quad Ax \le b , \end{array}$$

• If we perturb (A,b) to $(A+\Delta A,b+\Delta b)$ and z' is the new optimal value, then

$$|z' - z^*| \le C(d)^2 \| (\Delta A, \Delta b) \|$$

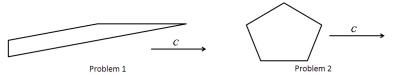
Sensitivity of solutions: The impact of geometry

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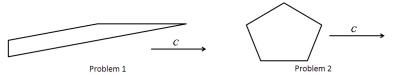
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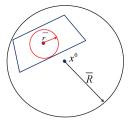
- problem 1 can be much more sensible (less robust) than problem 2 (for the current objective function)
- Intuition: geometric shape has something to say about sensitivity.
- Can we describe the way in which the optimal value of the problem changes with changes in the data, using geometric measures?

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- r is such that $B(\bar{x},r) \subset \mathcal{F}$, for some \bar{x} .



• The number R/r is an "aspect ratio" of the set S.

Changes in the objective function

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Theorem

Let $\Delta d = (\Delta A, \Delta b, 0)$ be a perturbation of the problem instance. Let z(d) and $z(d + \Delta d)$ be the corresponding optimal values. Then,

$$|z(d) - z(d + \Delta d)| \le 2\|c\| \left(\frac{R}{r}\right) \left(\frac{1}{\gamma(A)}\right) \times (\|\Delta b\| + \|\Delta A\|(R + \|x^0\|))$$

• where $\gamma(A)$ is a number depending on the matrix A.

The question: introducing those measures in the 2-stage problem

- Either C(d) or R/r could be used in connection to the function s we postulated before, but the dependence on the data has to be made explicit.
- Recall the conceptual problem:

$$TP_{R2}) \quad \min_{t=1}^{T} \sum_{t=1}^{T} C_t(\omega_t, x_t, y_t) + \sum_{t=1}^{T} s(\bar{\omega}_t, x_t, y_t)$$

s.t. $G_t(\omega_t, x_t, y_t) \le b_t$ $t = 1, ..., T$
 $H(\omega, x, y) = 0$

- The structure of this problem will be complicated, but it could be possible to construct bounds that limit the variation of *s* in terms of the tactical decisions.
- So far, this is an open problem, at least for me...

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• So far, this is also an open problem.

Final Coments

- We have been exploring consistency issues in the use of Optimization models in an intertemporal setting.
- Robust Optimization and the classical 2-stage stochastic approach have been promising on helping to achieve consistency.
- We plan to work on the general formulation considering direct consideration of stability measures.
- Our current "test bed" is a production planning problem in the forest industry.
- We are beginning to work in a second problem related to planning capacity in an hospital, where there are many sources un uncertainty.

THANKS!!

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