# Recent advances on the acceleration of first-order methods in convex optimization

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- Basic first-order descent methods
- Nesterov's acceleration
- Dynamic interpretation
  - Damped Inertial Gradient System (DIGS)
- Properties of DIGS trajectories and accelerated algorithms
- A first-order variant bearing second-order information in time and space

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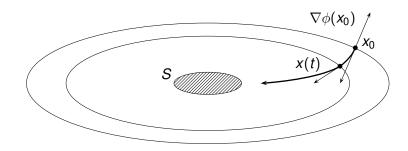
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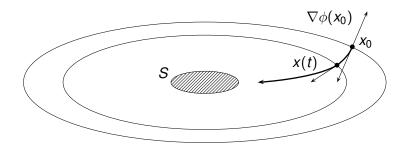
# **BASIC DESCENT METHODS**

Steepest descent dynamics:  $\dot{x}(t) = -\nabla \phi(x(t)), x(0) = x_0$ 



$$\frac{d}{dt}\phi(x(t)) = \langle \nabla \phi(x(t)), \dot{x}(t) \rangle = -\|\nabla \phi(x(t))\|^2 = -\|\dot{x}(t)\|^2$$

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Explicit discretization → gradient method (Cauchy 1847):

$$\frac{x_{k+1}-x_k}{\lambda}=-\nabla\phi(x_k)\quad\Longleftrightarrow\quad x_{k+1}=x_k-\lambda\nabla\phi(x_k).$$

Implicit discretization  $\rightarrow$  proximal method (Martinet 1970):

$$\frac{z_{k+1}-z_k}{\lambda}=-\nabla\phi(z_{k+1})\quad\Longleftrightarrow\quad z_{k+1}+\lambda\nabla\phi(z_{k+1})=z_k.$$

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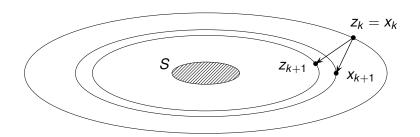
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#### Gradient

$$x_{k+1} = x_k - \lambda \nabla \phi(x_k)$$

#### Proximal

$$z_{k+1} + \lambda \nabla \phi(z_{k+1}) = z_k$$

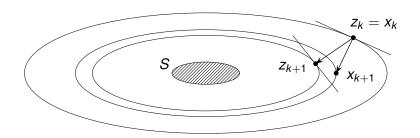


#### Gradient

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### Pros and cons

#### Gradient method

- Lower computational cost per iteration (explicit formula), easy implementation
- Convergence depends strongly on the regularity of the function (typically  $\phi \in \mathcal{C}^{1,1}$ ) and on the step sizes

#### Proximal point algorithm

- + More stability, convergence certificate for a larger class of functions  $(\nabla \phi \to \partial \phi)$ , independent of the step size
- Higher computational cost per iteration (implicit formula), often requires inexact computation

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#### Problem

$$\min\{\Phi(x):=F(x)+G(x):x\in H\},$$

where F is not smooth but G is.

Forward-Backward Method (
$$x_k o x_{k+rac{1}{2}} o x_{k+1}$$
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$$x_{k+1} + \lambda \partial F(x_{k+1}) \ni x_{k+\frac{1}{2}} = x_k - \lambda \nabla G(x_k)$$

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Gradient projection:

Goldstein 1964, Levitin-Polyak 1966, with  $F = \delta_C$ 

General setting:

Lions-Mercier 1979, Passty 1979

Iterative Shrinkage-Thresholding Algorithm (ISTA): Daubechies-Defrise-DeMol 2004, Combettes-Wajs 2005, for " $\ell^1 + \ell^2$ " minimization

$$\Phi(x) = F(x) + G(x) = \mu ||x||_1 + \frac{1}{2} ||Ax - b||^2$$

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# Convergence of the forward-backward method

#### Theorem

Let  $\Phi = F + G$ , where G is closed and convex, and F is convex with  $\nabla F$  L-Lipschitz. Assume  $\Phi$  has minimizers, and let  $(x_k)$  be obtained by the FB method with  $\lambda \leq 1/L$ . Then

- As  $k \to \infty$ ,  $(x_k)$  converges\* to a minimizer of  $\Phi$ ; and
- $\Phi(x_k)$  min  $\Phi = \mathcal{O}(k^{-1})$ : There is C > 0 such that

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# Convergence ISTA

Let  $\Phi : \mathbb{R}^N \to \mathbb{R}$  be defined by

$$\Phi(x) = \|x\|_1 + \frac{1}{2}\|Ax - b\|^2.$$

Local linear convergence results have been found recently, as well as theoretical convergence rates.

#### Theorem (Bolte-Nguyen-P.-Suter 2015

Let  $(x_k)$  be obtained by the FB method with step size  $\lambda$ . Then, there is an explicit constant d such that

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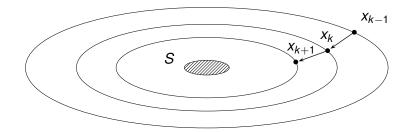
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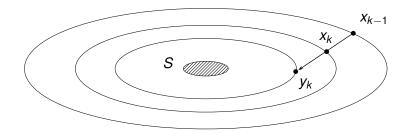
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# **NESTEROV'S ACCELERATION**

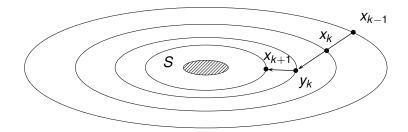
The main idea is the following: Instead of doing



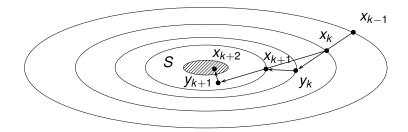
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- Convergence and its rate are sensitive to the choice of  $y_k$
- This simple procedure (Nesterov 1983) can take the theoretical rate of worst-case convergence for the values from the typical O(1/k) down to  $O(1/k^2)$
- No convergence proof for the iterates x<sub>k</sub>
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### ISTA & FISTA

#### General case:

- FB: values  $\mathcal{O}(k^{-1})$ , convergent sequence.
- AFB: values  $\mathcal{O}(k^{-2})$ .

#### $\ell^1 + \ell^2$ minimization:

- ISTA: values  $\mathcal{O}(Q^k)$ , convergent sequence (proved).
- FISTA: values (observed, not proved) \( \mathcal{O}(\tilde{Q}^k), \) always strictly faster than ISTA, convergent sequence (observed, not proved).

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- Is  $\Phi(x_k) \min \Phi = \mathcal{O}(\tilde{Q}^k)$  true for FISTA  $(\ell^1 + \ell^2)$ ?
- Is AFB always strictly faster than FB?
- What about FISTA and ISTA?
- Is  $\Phi(x_k)$  min  $\Phi = \mathcal{O}(k^{-2})$  optimal for AFB (in general)?
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# **DYNAMIC INTERPRETATION**

A finite-difference discretization of

(DIGS) 
$$\ddot{x}(t) + \frac{\alpha}{t}\dot{x}(t) + \partial F(x(t)) + \nabla G(x(t)) \ni 0.$$

gives

$$\frac{1}{h^2}(x_{k+1}-2x_k+x_{k-1})+\frac{\alpha}{kh^2}(x_k-x_{k-1})+\partial F(x_{k+1})+\nabla G(y_k)\ni 0,$$

where  $y_k$  (specified later) is related to the segment  $[x_{k-1}, x_k]$ .

#### Rewriting

$$\frac{1}{h^2}(x_{k+1}-2x_k+x_{k-1})+\frac{\alpha}{kh^2}(x_k-x_{k-1})+\partial F(x_{k+1})+\nabla G(y_k)\ni 0,$$

with  $\lambda = h^2$ , we obtain

$$x_{k+1} + \lambda \partial F(x_{k+1}) \ni x_k + \left(1 - \frac{\alpha}{k}\right)(x_k - x_{k-1}) - \lambda \nabla G(y_k).$$

Thus, if we set  $y_k = x_k + \left(1 - \frac{\alpha}{k}\right)(x_k - x_{k-1})$ , we obtain

$$X_{k+1} + \lambda \partial F(X_{k+1}) \ni Y_k - \lambda \nabla G(Y_k)$$

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Therefore, a finite-difference discretization of

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naturally yields

$$\begin{cases} y_k = x_k + \left(1 - \frac{\alpha}{k}\right)(x_k - x_{k-1}) \\ x_{k+1} = \operatorname{Prox}_{\lambda F} \circ \operatorname{Grad}_{\lambda G}(y_k) \end{cases}$$

Construction due to Su-Boyd-Candès 2014

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### PROPERTIES OF DIGS TRAJECTORIES

### Basic properties

#### Theorem (Attouch-Chbani-P.-Redont 2015)

If  $\alpha > 0$ , then

- $\bullet \lim_{t\to +\infty} \Phi(x(t)) = \inf(\Phi) \in \mathbb{R} \cup \{-\infty\}.$
- Every weak limit point of x(t), as  $t \to \infty$ , minimizes  $\Phi$ .
- Either  $\Phi$  has minimizers and all trajectories are bounded, or it does not and all trajectories diverge to  $+\infty$  in norm.
- If  $\Phi$  is bounded from below, then  $\lim_{t\to +\infty} \|\dot{x}(t)\| = 0$ .

# Rate of convergence

#### Theorem (Su-Boyd-Candès 2014)

If  $\alpha \geq 3$  and  $\Phi$  has minimizers, then every solution satisfies

$$\Phi(x(t))-\min(\Phi)\leq \frac{C}{t^2},$$

where C depends on  $\alpha$  and the initial data.

# Rate of convergence

The exponent 2 is sharp. More precisely, we have the following:

#### Theorem (ACPR)

For each p>2, there is  $\Phi$  such that  $\Phi$  has minimizers and every solution satisfies

$$\Phi(x(t)) - \min(\Phi) = \frac{C}{t^p}.$$

### Rate of convergence

If  $\Phi$  is strongly convex, convergence is arbitrarily fast, as  $\alpha$  grows.

#### Theorem (ACPR)

Let  $\Phi$  be strongly convex and let  $x^*$  be its unique minimizer. Every solution satisfies

$$\Phi(x(t)) - \min(\Phi) \leq \frac{C}{t^{\frac{2}{3}\alpha}} \quad \text{and} \quad \|x(t) - x^*\| \leq \frac{D}{t^{\frac{1}{3}\alpha}},$$

where C and D depend on  $\alpha$ , the strong convexity parameter and the initial data.

#### Theorem (ACPR, May)

- x(t) converges weakly, as  $t \to +\infty$ , to a minimizer of  $\Phi$ .
- Convergence is strong if either Φ is uniformly convex, int(Argmin(Φ)) ≠ ∅, or Φ is even.
- $\|\dot{x}(t)\| = o(t^{-1}).$
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# PROPERTIES OF ACCELERATED ALGORITHMS

#### Recall that

$$\begin{cases} y_k = x_k + \left(1 - \frac{\alpha}{k}\right)(x_k - x_{k-1}) \\ x_{k+1} = \operatorname{Prox}_{\lambda F} \circ \operatorname{Grad}_{\lambda G}(y_k) \end{cases}$$

#### Theorem (ACPR)

If  $\alpha > 0$ , then

- $\lim_{k\to+\infty}\Phi(x_k)=\inf(\Phi)$ ; and
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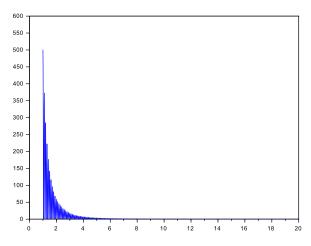
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We consider the function  $\Phi(x_1, x_2) = \frac{1}{2}(x_1^2 + 1000x_2^2)$ . We show the behavior of a solution to

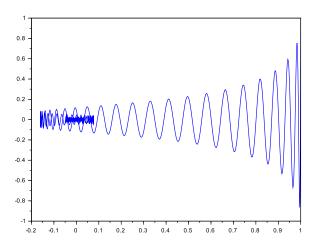
$$\ddot{x}(t) + \frac{\alpha}{t}\dot{x}(t) + \nabla\Phi(x(t)) = 0$$

on the interval [1, 20] with  $\alpha =$  3.1 .

### **Function values**



# Trajectory



# CAN WE DO BETTER?

# Idea: Newton / Levenberg-Marquardt

#### Pros:

- Is fast.
- Compensates the effect of ill-conditioning.

#### Cons

- Requires higher regularity (to compute and invert the Hessian).
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## **NDIGS**

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Seems much more complicated, but

#### Proposition (APR 2015)

System (NDIGS) is equivalent to

$$\begin{cases} \dot{x}(t) + \beta \nabla \Phi(x(t)) - \left(\frac{1}{\beta} - \frac{\alpha}{t}\right) x(t) + \frac{1}{\beta} y(t) = 0 \\ \dot{y}(t) - \left(\frac{1}{\beta} - \frac{\alpha}{t} + \frac{\alpha\beta}{t^2}\right) x(t) + \frac{1}{\beta} y(t) = 0 \end{cases}$$

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Using variable Z = (x, y), this is

$$\dot{Z}(t) + \nabla \mathcal{G}(Z(t)) + D(t, Z(t)) \ni 0,$$

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#### Theorem (APR)

Let  $\Phi$  be closed and convex, and let  $\beta > 0$ .

- All the conclusions obtained for the solutions of (DIGS) are also true for the solutions of (NDIGS').
- But also  $\lim_{t\to\infty} \|\nabla \Phi(x(t))\| = 0$ .
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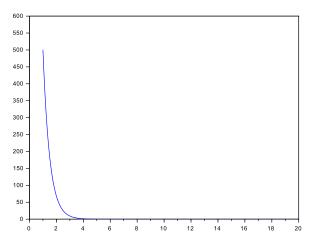
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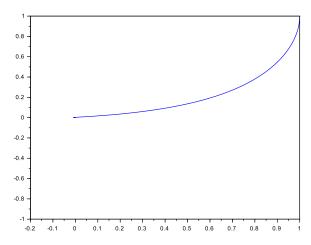
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# Algorithmic implementation

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